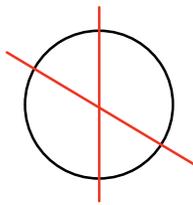
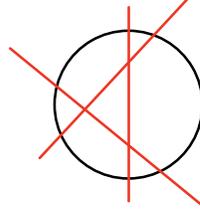


Cake cutting problem
 What is the greatest number of pieces that a cake can be cut into with a given number of cuts? ($d=2$)



2 cuts
 $P_2(2) = 4$

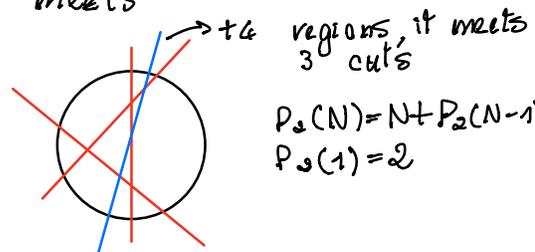
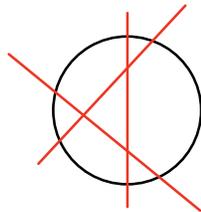


3 cuts
 $P_2(3) = 7 \dots$

Cuts N	1	2	3	4	5	...
Pieces $P_2(N)$	2	4	7	11	16	...

$$\begin{cases} 2 = 1 + 1 \\ 4 = 1 + 1 + 2 \\ 7 = 1 + 1 + 2 + 3 \\ \vdots \end{cases}$$

What it seems is that, having done $n-1$ "maximal" cuts, the number of pieces subdivided by the N -th cut equals one more than the number of previous cuts it meets



$$\begin{aligned} P_2(N) &= N + P_2(N-1) \\ P_2(1) &= 2 \end{aligned}$$

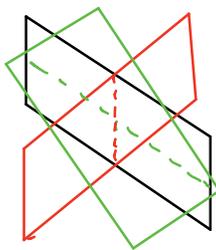
We can give a close formula to $P_2(N)$:

$$P_2(N) = 1 + \sum_{i=1}^N 1 = 1 + \frac{N(N+1)}{2}$$

What about general d ?

What we're asking is to compute the number of cells in an arrangement of n hyperplanes in d -dimensions in general position.

We can find a recursion relation like before, but now $d-1$ dimensional hyperplanes intersect and form $d-2$ dimensional objects, which were points in $d=2$.



The N -th hyperplane will intersect $n-1$ previous ones and creates $P_{d-1}(N-1)$ new cells:

$\cdot d=2$ $P_1(N-1)$ $= N$

$\cdot d=3$ $P_2(N-1)$ $= \binom{N-1}{0} + \dots + \binom{N-1}{2}$

Check: $N=3 \Rightarrow \binom{2}{0} + \binom{2}{1} + \binom{2}{2} = 1 + 2 + 1 = 4$

In essence, we have $P_d(N) = P_d(N-1) + P_{d-1}(N-1)$

Claim: $P_d(N) = \binom{N}{0} + \binom{N}{1} + \dots + \binom{N}{d}$

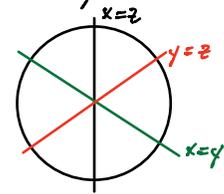
Proof: $\binom{N}{0} + \binom{N}{1} + \dots + \binom{N}{d-1} + \binom{N}{d} \stackrel{?}{=} \binom{N-1}{0} + \binom{N-1}{1} + \dots + \binom{N-1}{d-1} + \binom{N-1}{d} \quad \left. \vphantom{\binom{N-1}{0}} \right\} P_d(N-1)$
 $+ \binom{N-1}{0} + \binom{N-1}{1} + \dots + \binom{N-1}{d-1} \quad \left. \vphantom{\binom{N-1}{0}} \right\} P_{d-1}(N-1)$

We just need to use $\binom{N-1}{k} + \binom{N-1}{k+1} = \binom{N}{k+1}$ □

Symmetric cake-cutting

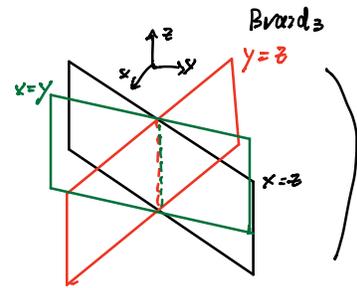
What are the possible ways to cut a perfectly round cake so that all pieces are congruent?

This can be reformulated as a problem of hyperplane arrangements, and in particular braid arrangement in the space of dimension $d = \text{number of cuts}$



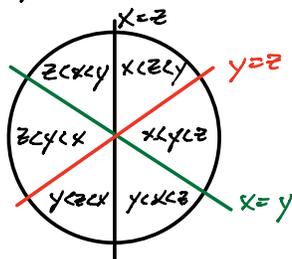
$d = \text{number of cuts}$

(i.e. you can see the figure as a $d=2$ projection of

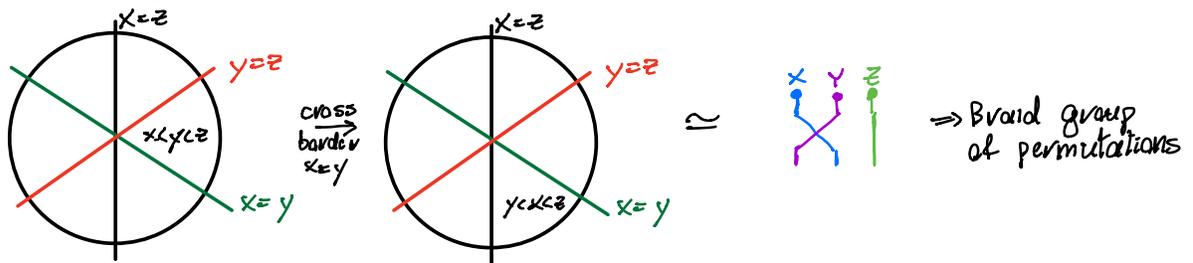


Each region of $Braid_{d=3}$ is on one side of the planes $x=y, x=z, y=z$. Therefore

- either $x < y$ or $y < x$
 - either $x < z$ or $z < x$
 - either $y < z$ or $z < y$
- \Rightarrow 6 possibilities
- | | | |
|-------------|-------------|-------------|
| $x < y < z$ | $y < x < z$ | $z < x < y$ |
| $x < z < y$ | $y < z < x$ | $z < y < x$ |



Note: crossing one border means flipping one inequality



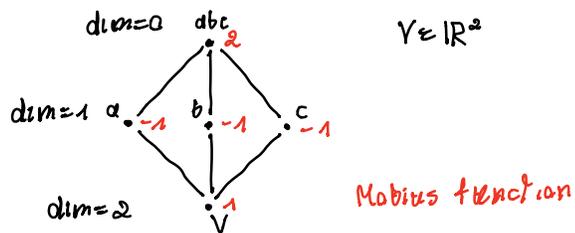
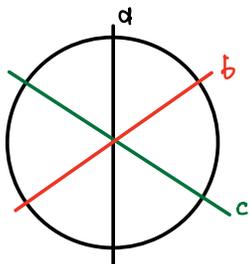
How many regions?

$$\square < \square < \square$$

There are 3 possibilities for the first letter in the inequality, two for the second one and one for the last one

Sol. **6** ($= 3!$)

This counting problem for arrangement has a nice formulation in terms of the characteristic polynomial. This will be a subject of the lecture next week, but since there will be no tutorial next week I thought it would be good to make a simple example here. Let's take the intersection poset of the arrangement ordered by reverse inclusion



The characteristic polynomial of an arrangement is Def. (Stanley): $\chi_A(t) = \sum_{x \in L(A)} \mu(x) t^{\dim(x)}$

Let's apply it to our case:

$$1 \cdot t^2 - 3 \cdot t + 2 = \chi_A(t) = t^2 - 3t + 2$$

Theorem (Zaslavsky, 1975): $|\chi_A(-1)| =$ number of regions formed by the arrangement A in \mathbb{R}^d

$$|\chi_A(-1)| = \mathbf{6} \quad \checkmark$$

Even more

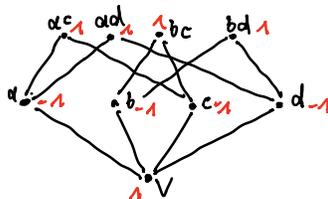
Theorem (Zaslavsky): $|\chi_A(t)|$ = number of bounded regions formed by the arrangement A in \mathbb{R}^d

In the previous case

$$|\chi_A(t)| = 0 \quad \checkmark$$

Let's consider a non-trivial example

	a	b	
c	1	2	3
	ac	bc	
d	4	5	6
	ad	bd	
	7	8	9



$$\chi_A(t) = t^2 - 4t + 4$$

$$|\chi_A(+1)| = 1 \quad \checkmark$$

$$|\chi_A(-1)| = 9 \quad \checkmark$$